## NAME:

## STUDENT NUMBER:

## Question 1, State Space and Heuristics, (10 marks)

In the context of the 8-puzzle problem give a description and a data structure for the following:
i) State space representation.
ii) Legal move representation.

## Solution

i) We could represent any state of the 8 -puzzle by a permutation of the nine integers, $[0,1,2,3,4,5,6,7,8]$. The position of the 0 integer in the permutation indicates the open slot in the puzzle.
For example: The start and goal states

| 2 | 8 | 3 |
| :--- | :--- | :--- |
| 1 | 6 | 4 |
| 7 |  | 5 |


| 1 | 2 | 3 |
| :--- | :--- | :--- |
| 8 |  | 4 |
| 7 | 6 | 5 |

could be represented by the permutations
$[2,8,3,1,6,4,7,0,5]$ and $[1,2,3,8,0,4,7,6,5]$
Full state space is then the tree of reachable states starting from the initial state and with each state connected by a legal move. It is not necessary to store the full state space in order to solve the 8-puzzle, it is only necessary to search state space for the goal state.
ii) A move for the 8-puzzle is represented by moving the empty slot in one of 4 directions, $[\mathrm{U}, \mathrm{D}, \mathrm{L}, \mathrm{R}]$, The move is legal if the direction selected does not move the empty slot off the $3 \times 3$ grid.

## Question 2, Search Algorithms, (10 marks)

Using Best First Search, construct an $A^{*}$ algorithm for the 8-puzzle problem and explain why your algorithm is indeed $A^{*}$.

## Solution

```
def bestfs(start,goal)
    open = [start],
    close = [],
    State = failure;
    while (open <> []) AND (State <> success) begin
        remove the leftmost state from open, call it X;
        if X is the goal, then
            State = success
        else begin
            generate children of X;
            for each child of X
                do case
                        the child is not on open or closed
                            assign the child a heuristic value,
                            add the child to open,
                    the child is already on open
                    if the child was reached by a shorter path then
                                give the state on open the shorter path
                    the child is already on closed:
                    if the child is reached by a shorter path then
                        remove the state from closed and add the child to open;
                endcase
            endfor
            put X on closed;
        end;
        re-order states on open by (path_length_so_far + heuristic) merit (best leftmost);
    endwhile;
    return State;
end.
def heuristic(state,goal)
    return(number_tiles_out_of_place(state,goal))
end.
```

This algorithm is $A^{*}$ because the next state searched has minimum path-length-so-far + heuristic from here-to-goal and the heuristic under-estimates the actual number of moves required to take the current state to the goal state.

## Question 3, Game Playing, (10 marks)

Consider Grundy's two player version of NIM starting with 8 -counters in one pile.
a) Draw the full game tree.
b) Use the minimax search algorithm to show that max can always win if he plays first.

## Solution

See notes......

## Question 4, Prolog, (10 marks)

Consider the following first verse of a well-known hillbilly poem by Moe Jaffe:
Many many years ago, when I was twenty-three, I married the widow, Hoe, as pretty as can be,
The widow had a grown-up daughter, Roe, with hair of red, My father, Poe, fell for her and soon the two were wed.
a) Construct Prolog predicates to capture the relationships in the poem.
b) Construct a grandfather predicate in the loose sense that allows for step-relationships.
c) Explain how you would use your grandfather predicate to prove that Moe was now his own grandfather.

## Solution

```
male(moe).
male(poe).
female(hoe).
female(roe).
spouse(moe, hoe).
spouse (hoe,moe).
spouse (poe,roe).
spouse (roe,poe).
parent(moe,poe).
parent(roe,hoe).
stepParent(X,Y) :-
    parent (X,W),
    spouse(W,Y).
grandfather(X,Y) :-
    male(Y),
    (parent (X,W); stepParent (X,W)),
    (parent(W,Y); stepParent (W,Y)).
```


## Question 5, Logic, (10 marks)

Convert the following statements into WFFs.

1) Animals, that do not kick, are always unexcitable
2) Donkeys have no horns
3) A buffalo can always toss one over a gate
4) No animals that kick are easy to swallow
5) No hornless animal can toss one over a gate
6) All animals are excitable, except buffalo
and the use resolution to show that

Donkeys are not easy to swallow.

## Solution

| 1) | $\forall x: \sim \operatorname{kicker}(x) \rightarrow \sim \operatorname{excitable}(x)$ | kicker $\left(a_{1}\right) \vee \sim \operatorname{excitable}\left(a_{1}\right)$ |
| :---: | :---: | :---: |
| 2) | $\forall x: \operatorname{donkey}(x) \rightarrow \sim \operatorname{horned}(x)$ | $\sim \operatorname{donkey}\left(a_{2}\right) \vee \sim \operatorname{horned}\left(a_{2}\right)$ |
| 3) | $\forall x: \operatorname{buffalo}(x) \rightarrow \operatorname{tosser}(x)$ | $\sim \operatorname{buffalo}\left(a_{3}\right) \vee$ tosser $\left(a_{3}\right)$ |
| 4) | $\forall x: \operatorname{kicker}(x) \rightarrow \sim \operatorname{swallow}(x)$ | $\sim \operatorname{kicker}\left(a_{4}\right) \vee \sim \operatorname{swallow}\left(a_{4}\right)$ |
| 5) | $\forall x: \sim \operatorname{horned}(x) \rightarrow \sim \operatorname{tosser}(x)$ | horned ( $\left.a_{5}\right) \vee \sim \operatorname{tosser}\left(a_{5}\right)$ |
| 6) | $\forall x: \operatorname{excitable}(x) \rightarrow \sim \operatorname{buffalo}(x)$ | $\sim \operatorname{excitable}\left(a_{6}\right) \vee \sim \operatorname{buffalo}\left(a_{6}\right)$ |
| 7) | $\forall x: \sim \operatorname{buffalo}(x) \rightarrow$ excitable $(x)$ | buffalo $\left(a_{7}\right) \vee$ excitable $\left(a_{7}\right)$ |
|  | negate the goal |  |
| 8) | $\sim(\forall x: \operatorname{donkey}(x) \rightarrow \sim \operatorname{swallow}(x))$ | donkey(a) |
| 9) |  | swallow(a) |
|  | resolution |  |
| 10) | 8 and 2 | $\sim \operatorname{horned}(a)$ |
| 11) | 10 and 5 | $\sim \operatorname{tosser}(\mathrm{a})$ |
| 12) | 11 and 3 | $\sim$ buffalo(a) |
| 13) | 12 and 7 | excitable(a) |
| 14) | 9 and 4 | $\sim \operatorname{kicker}(a)$ |
| 15) | 14 and 1 | $\sim$ excitable (a) |
| 16) | 15 and 13 | FALSE |

## Attachments

Summary of first order logic algebra:

$$
\begin{aligned}
& A \rightarrow B \equiv(\sim A) \vee B \\
& \sim(\sim A) \equiv A \\
& \sim(A \wedge B) \equiv(\sim A) \vee(\sim B) \\
& \sim(A \vee B) \equiv(\sim A) \wedge(\sim B) \\
& \sim(\forall x: P(x)) \equiv \exists x: \sim P(x) \\
& \sim(\exists x: P(x)) \equiv \forall x: \sim P(x)
\end{aligned}
$$

